

Validation of Abstract Side-channel Models for Computer Architectures

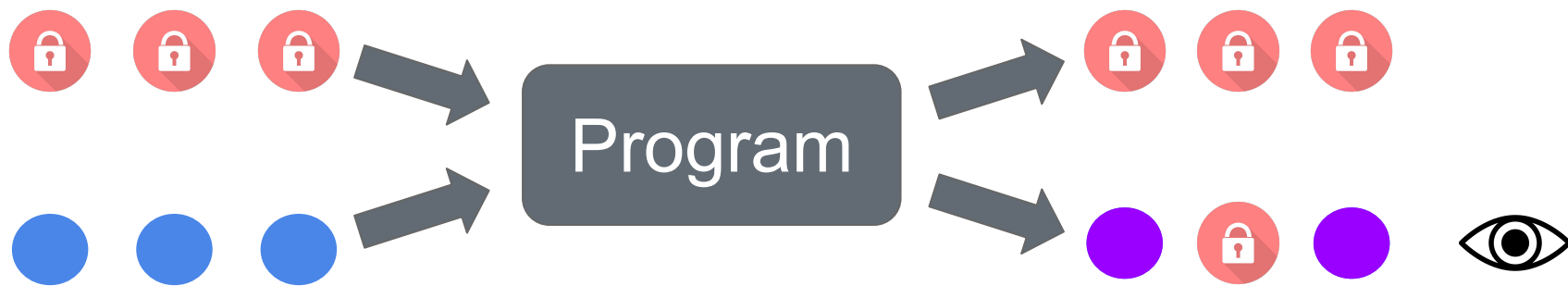
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(work in progress)
ENTROPY '19,
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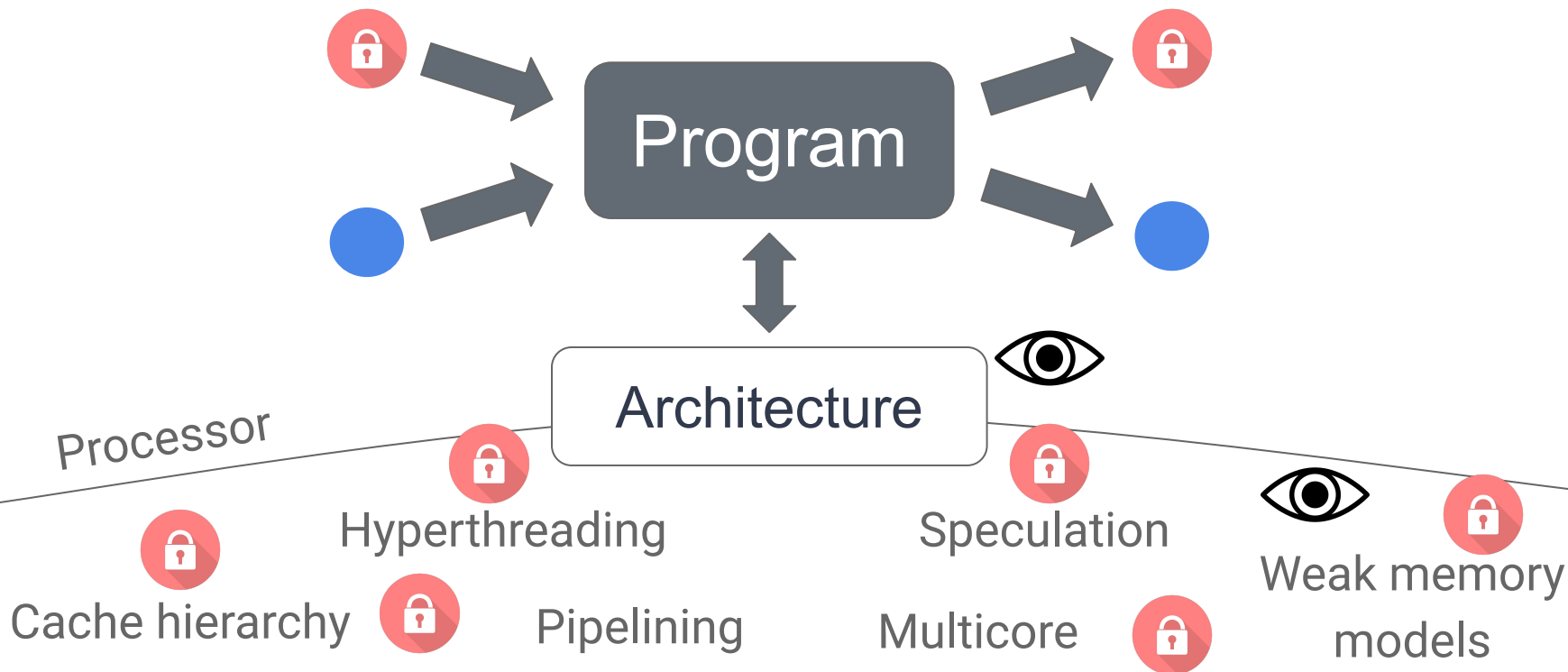
Side channels and Security



Noninterference

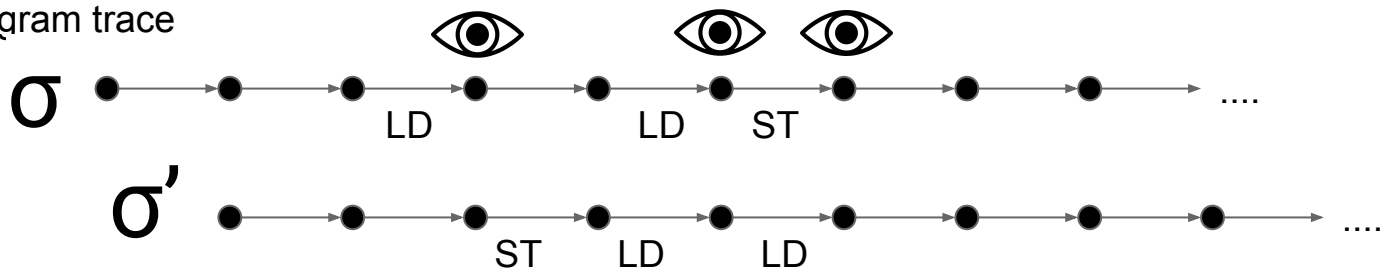
$P(a) \sim_{\text{eye}} P(b)$ if $a \sim_{\text{eye}} b$, $\forall a b$

Side channels and Security



Abstract observation models

Program trace



Program text

MOV ...
ADD ...
CMP ...
LD ...
B label
...



$\sigma \sim_P \sigma'$

Observational
Equivalence

Abstract model

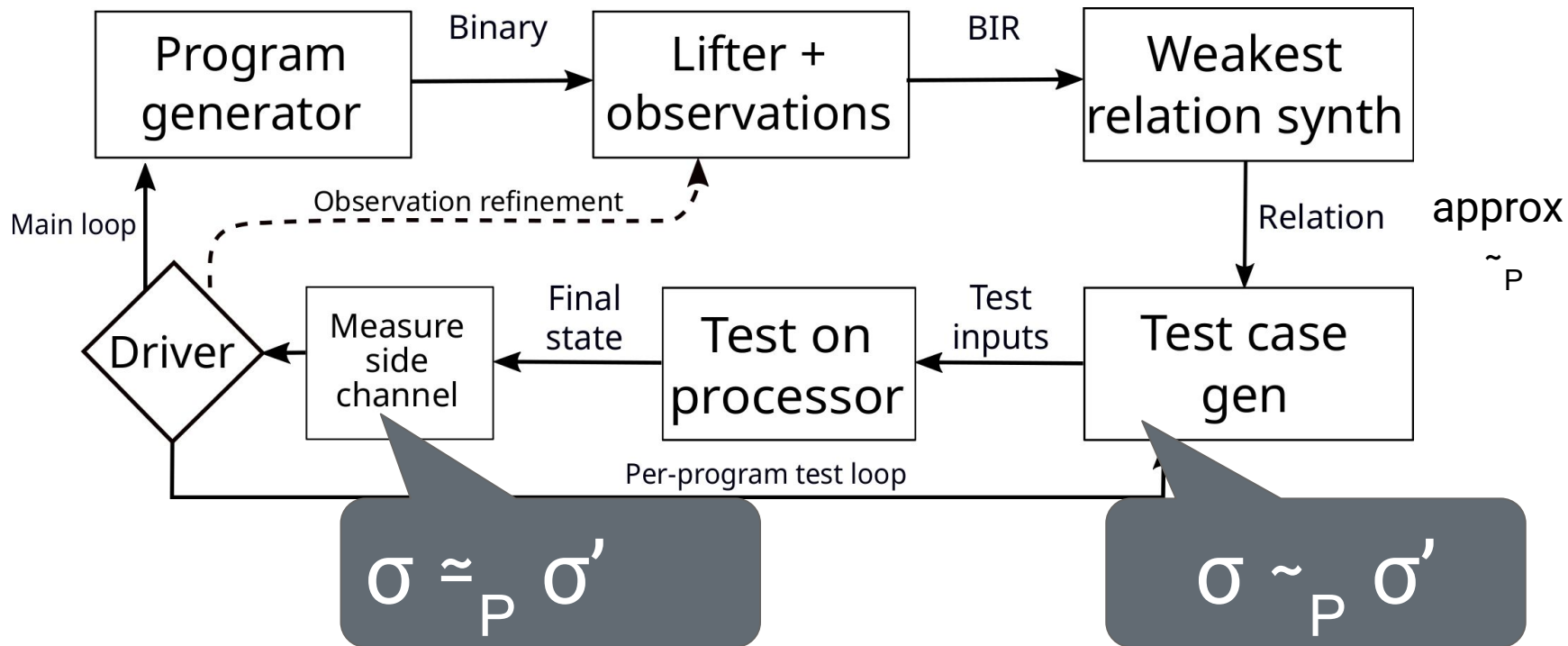
Real hardware

$$\sigma \sim_P \sigma' \quad \Rightarrow \quad \sigma \approx_P \sigma'$$

?

**SCAM-V: Side Channel
Abstract Model Validator**

SCAM-V Pipeline Overview



BIR

Abstract Assembly Language

- Infinite number of register variables
- Assignments, jumps, cond. jumps
- BIR expressions: arithmetic, bitwise, etc.
- Memory is an array
- **(Attacker) Observations**

BIR

Observation statements

$\text{OBS}(c, \text{exp})$

outputs the evaluation
of exp in the current
state,
if c evaluates to true

$\text{OBS}(\text{true}, \text{exp}) = \text{OBS}(\text{exp})$

Lifting and BIR

Binary

```
B.eq 12  
mul x1 x2 x3  
12: ldr x2 {x1} +8
```

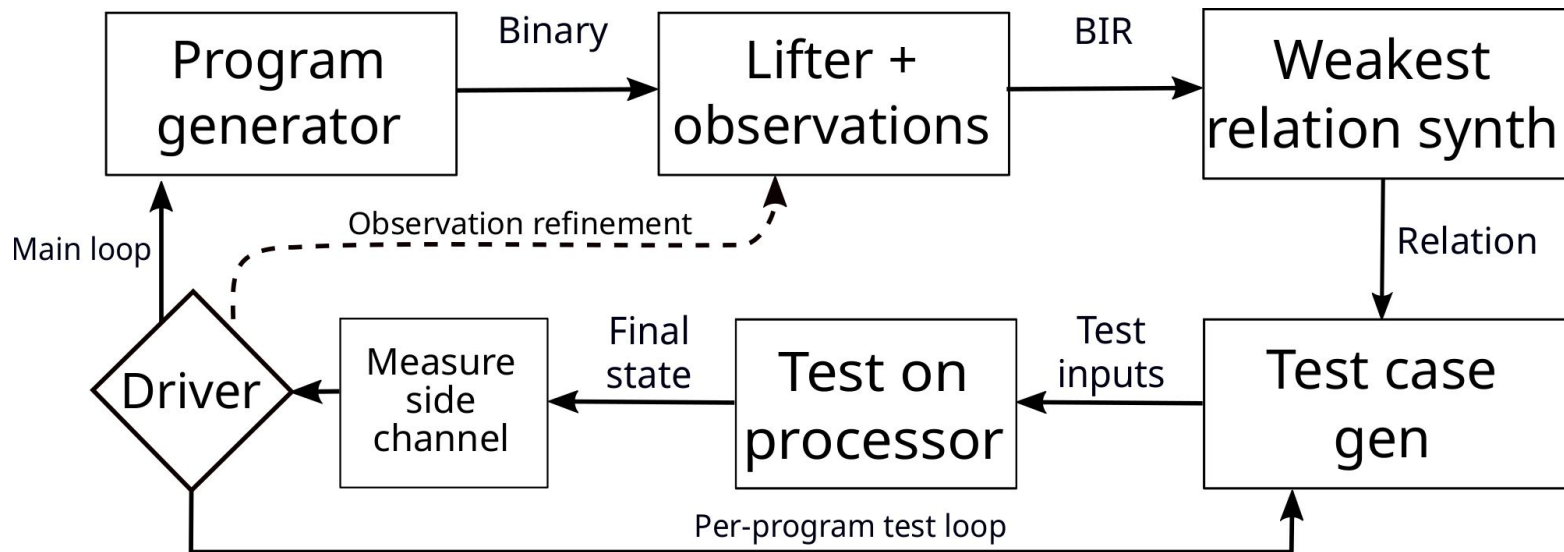
BIR

```
————→ [10: CJMP Z 12 11]  
————→ [11: X1= X2*X3; JMP 12]  
————→ [12: OBS([tag(X1)]);  
X2= LOAD MEM, X1); X1= X1+8;  
HALT]
```

Obs model:
Cache tag

Observation added
according to
selected model

SCAM-V Pipeline Overview



SCAM-V Weakest Relation Synthesis

```
[10: CJMP Z 12 11]
[11: X1= X2*X3; JMP 12]
[12: OBS([tag(X1)]);
      X2= LOAD(MEM, X1); X1= X1+8;
      HALT]
```



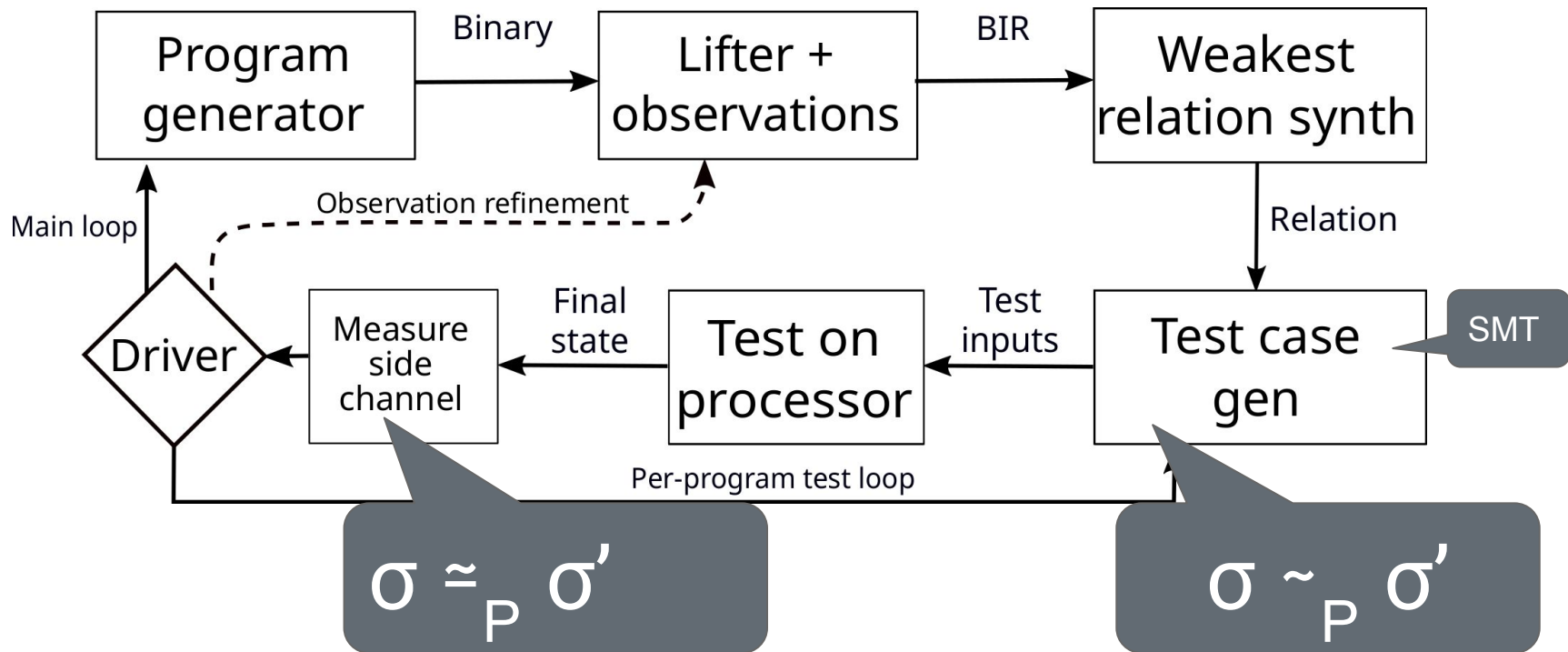
Symbolic execution



Relation

$$\sigma \sim_P \sigma'$$
$$\begin{aligned} & (Z \wedge Z' \Rightarrow \text{tag}(X1) = \text{tag}(X1')) \\ & \wedge (Z \wedge \neg Z' \Rightarrow \text{tag}(X1) = \text{tag}(X2'*X3')) \\ & \wedge (\neg Z \wedge Z' \Rightarrow \text{tag}(X2*X3) = \text{tag}(X1')) \\ & \wedge (\neg Z \wedge \neg Z' \Rightarrow \text{tag}(X2*X3) = \text{tag}(X2'*X3')) \\ & \wedge \text{address constraints } (\dots) \end{aligned}$$

SCAM-V Pipeline Overview



Testing / Measuring the channel

- ARMv8 (Raspberry Pi 3)
- **Observation:** tag and operation of every memory access
- **Measuring:** TrustZone instructions for cache inspection
- Cortex-M0 (MicroBit)
- **Observation:** program counter
- **Measuring:** internal system clock

SCAM-V: Side Channel Abstract Model Validator

Summary

- Modern architectures are too complex to directly analyze side-channels
- Abstract models based on system-state observations
 - PC security model, Cache-line + tag of memory access
- **Assumption:** States with **equivalent** observations in the model are **indistinguishable** to the attacker on real hardware
- Not always true! e.g. Spectre
- SCAM-V **validates** this assumption